

Log-Concave Polynomials I

Week 5: Reading Group W20

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Agenda

- Generating polynomial $g_M(z)$ is CLC
- CLC distributions and entropy
- Max. entropy programs and basis counting

I. Generating polynomial $g_M(z)$ is CLC

- Chow ring, graded Möbius algebra, and polynomials
- From $\nabla^2 g_M(z)$ to $\nabla^2(\log g_M(z))$

Known fact about Chow ring

- Given (simple) matroid M , on $E = [n]$, rank r
- $\mathbb{A}^*(M)$: Chow ring over the nonempty proper *flats* F of M
- It's an algebra (\approx vector space + ring)
- Elements $\sum_F c_F x_F$, $c_F \in \mathbb{R}$, modulo two relations
 - $x_{F_1} x_{F_2} = 0$ if F_1, F_2 incomparable
 - $\sum_{i \in F} x_F = \sum_{j \in F} x_F$ for $i, j \in E$
- Graded using (homogeneous) degree of polynomial

The $\text{deg}(\cdot)$ map

$$\alpha_i = \sum_{i \in F} \chi_F, \text{ then } \alpha_i = \alpha_j = \alpha$$

① For any flag $F_1 \subsetneq F_2 \subsetneq \dots \subsetneq F_k$,
 $\exists j \in [k]$
 if $\text{rk}(F_j) > j$, then

$$\chi_{F_1} \chi_{F_2} \dots \chi_{F_k} \alpha^{r-k} = 0$$

② If $\text{rk}(F_j) = j$ for all $j \in [k]$,
 then $\chi_{F_1} \chi_{F_2} \dots \chi_{F_k} \alpha^{r-k} = \alpha^r$.

$$\chi_{\{i\}} \alpha^{r-1} = \alpha^r$$

$$\left(\sum_{F \ni i} \chi_F \right) \alpha^{r-1}$$

$$\chi_{F_1} \dots \chi_{F_k} \alpha^{r-k} = \chi_{F_i}$$

$j \in F_k \setminus F_{k-1}$

$$\left[\begin{array}{l} \text{Flat } F, i \notin F, \\ \chi_F \cdot \alpha = \chi_F \cdot \left(\sum_{G \supseteq F, i \in G} \chi_G \right) \\ \text{① } j = r-1 \checkmark \end{array} \right]$$

Fact. There exists an isomorphism
 $\text{deg}: \mathbb{A}^r(M) \rightarrow \mathbb{R}$

s.t. $\text{deg}(x_{F_1} x_{F_2} \dots x_{F_r}) = 1$ for any
 (maximal) flag $F_1 \subset F_2 \subset \dots \subset F_r$

Elements $\sum_F c_F \chi_F, c_F \in \mathbb{R}$, modulo:

- $\chi_{F_1} \chi_{F_2} = 0$ if F_1, F_2 incomparable
- $\sum_{i \in F} \chi_F = \sum_{j \in F} \chi_F$ for $i, j \in E$

The $\text{deg}(\cdot)$ map

Fact. There exists an isomorphism
 $\text{deg}: \mathbb{A}^r(M) \rightarrow \mathbb{R}$

s.t. $\text{deg}(x_{F_1} x_{F_2} \dots x_{F_r}) = 1$ for any
 (maximal) flag $F_1 \subset F_2 \subset \dots \subset F_r$

$$j \in F_k \setminus F_{k-1}$$

$$x_{F_1} \dots x_{F_k} \alpha^{r-k} \stackrel{?}{=} \alpha^r$$

$$\alpha^r = x_{F_1} \dots x_{F_{k-1}} \alpha^{r-(k-1)}$$

$$= x_{F_1} \dots x_{F_{k-1}} \left(\sum_{F \supset j} x_F \right) \alpha^{r-k}$$

$$= x_{F_1} \dots x_{F_{k-1}} x_{F_k} \alpha^{r-k}$$

Graded Möbius algebra

- $\mathbb{B}^*(M)$: over flats F , modulo the relations
 - $y_{F_1} y_{F_2} - y_{F_1 \vee F_2}$, for $\text{rank}(F_1) + \text{rank}(F_2) = \text{rank}(F_1 \vee F_2)$
 - $y_{F_1} y_{F_2}$, for $\text{rank}(F_1) + \text{rank}(F_2) > \text{rank}(F_1 \vee F_2)$
- For any basis I of a flat F , $y_F = \prod_{i \in I} y_i$

Relating basis counting and Chow ring

- Turns out $\mathbb{B}^*(M)$ can be embedded into $\mathbb{A}^*(M_0)$!
 - $y_i \mapsto \beta_i := \sum_{i \in F, 0 \notin F} x_F$, extend naturally
- For basis B of M , $\prod_{i \in B} y_i = \mathbf{1}$, so $\deg(\prod_{i \in B} \beta_i) = 1$ (up to scaling)
- For non-basis B of size r , $\prod_{i \in B} y_i = 0$
- **Prop. 3.3.** $\sum_{i=1}^n v_i \beta_i$ corresponds to ∂_v
- **Theorem 3.4.** For any $V \in \mathbb{R}_{\geq 0}^{n \times k}$, $0 \leq k \leq r - 2$, the Hessian of $D_V g_M(z)|_{z=\lambda}$ has at most one positive eigenvalue.

From $\nabla^2 g_M(z)$ to $\nabla^2 (\log g_M(z))$

- Let M be a simple matroid.
- **Theorem 3.4.** For any $V \in \mathbb{R}_{\geq 0}^{n \times k}$, $0 \leq k \leq r - 2$, the Hessian of $D_V g_M(z)|_{z=\lambda}$ has at most one positive eigenvalue.

Goal:

- For any $V \in \mathbb{R}_{\geq 0}^{n \times k}$, $0 \leq k \leq r - 2$, the Hessian of $\log D_V g_M(z)|_{z=\lambda}$ is negative semidefinite.

Whiteboard

Goal:

- For any $V \in \mathbb{R}_{\geq 0}^{n \times k}$, $0 \leq k \leq r - 2$, the Hessian of $\log D_V g_M(z)|_{z=\lambda}$ is negative semidefinite.

$$h(z) = D_V g_M(z).$$

Known: $\nabla^2 h(z)$ has ≤ 1 positive-eigenvalue

$$\nabla^2 \log h(z) = \frac{1}{h(z)} \left[\underbrace{h(z)}_{\text{Hessian}} \nabla^2 h(z) - \underbrace{(\nabla h(z))}_{\text{Hessian}} (\nabla h(z))^T \right]$$

Euler's identities:

f : degree d homogeneous poly,

$$\begin{aligned} \textcircled{1} \quad \langle \nabla f(z), z \rangle &= d \cdot f(z) \\ \textcircled{2} \quad \nabla^2 f(z) \cdot z &= (d-1) \cdot \nabla f(z) \\ \textcircled{3} \quad \underline{z^T \nabla^2 f(z) \cdot z} &= d(d-1) \cdot f(z) \end{aligned}$$

$$z_1^{d_1} z_2^{d_2} \dots z_n^{d_n} \begin{pmatrix} z_1 \partial_1 (\dots) \\ \vdots \\ z_n \partial_n (\dots) \end{pmatrix} = d_i (\dots)$$

Whiteboard

Goal:

- For any $V \in \mathbb{R}_{\geq 0}^{n \times k}$, $0 \leq k \leq r - 2$, the Hessian of $\log D_V g_M(z)|_{z=\lambda}$ is negative semidefinite.

$$h(z) \nabla^2 h(z) - (\nabla h(z)) (\nabla h(z))^T \quad (*)$$

Evaluated at $z = \lambda \in \mathbb{R}_{> 0}^n$.

$$\text{Set } A = \nabla^2 h(z)|_{z=\lambda}, \quad v = \lambda,$$

Then $(*)|_{z=\lambda}$

$$= \frac{1}{d(d-1)} (v^T A v) \cdot A - \frac{1}{(d-1)^2} (Av) (Av)^T.$$

Goal = Show that this is $\succeq 0$.

$$(v^T A v) \cdot A - t \cdot (Av) (Av)^T \succeq 0 \text{ for } t \geq 1.$$

Whiteboard

Goal:

- For any $V \in \mathbb{R}_{\geq 0}^{n \times k}$, $0 \leq k \leq r - 2$, the Hessian of $\log D_V g_M(z)|_{z=\lambda}$ is negative semidefinite.

Consider $u, v \in \mathbb{R}^n$.

$$u^T \left[(v^T A v) \cdot A - (A v)(A v)^T \right] u \leq 0.$$

$$\Leftrightarrow (v^T A v) \cdot (u^T A u) - (u^T A v)^2 \leq 0.$$

$$\Leftrightarrow \det \left(\begin{bmatrix} u^T \\ v^T \end{bmatrix} A \begin{bmatrix} u & v \end{bmatrix} \right) \leq 0$$

$$\begin{bmatrix} u^T A u & u^T A v \\ u^T A v & v^T A v \end{bmatrix}$$

① It has ≤ 1 positive e-value.

② It has ≥ 1 non-negative e-values.

(Some) affine maps preserve CLC

- Let $A \in \mathbb{R}_{\geq 0}^{n \times m}$, $b \in \mathbb{R}_{\geq 0}^n$. If $g(z_1, \dots, z_n) \in \mathbb{R}[z_1, \dots, z_n]$ is CLC, then $g(Az + b) \in \mathbb{R}[y_1, \dots, y_m]$ is also CLC.

$$g \circ T \quad Tz = Az + b$$

$$\begin{aligned} \textcircled{1} \quad g(T(\lambda z_1 + (1-\lambda)z_2)) &= g(\lambda T(z_1) + (1-\lambda)T(z_2)) \\ &\geq g(T(z_1))^\lambda \cdot g(T(z_2))^{1-\lambda}. \end{aligned}$$

$$\begin{aligned} \textcircled{2} \quad D_{\nu} (g(T(z))) \Big|_{z=\lambda} &= D_{A \cdot \nu} (g(y)) \Big|_{y=T(\lambda)} \\ &(\nu \geq 0) \end{aligned}$$

From simple matroids to general matroids

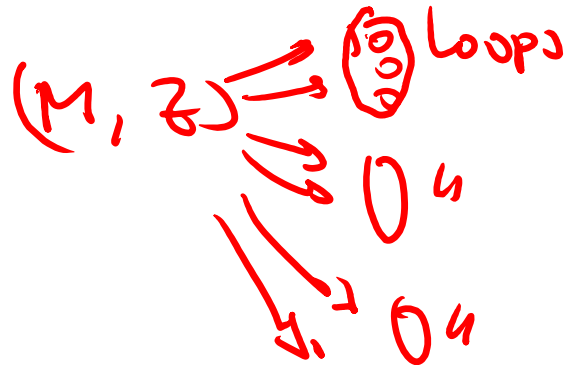
M : matroid (not necessarily simple)

~ simplify

\hat{M} : simple matroid

Delete loops, identify n elements

$g_{\hat{M}}$ is CCC. find T s.t. $g_M = g_{\hat{M}} \circ T$.



T maps loops to 0 .
maps parallel elements
to the representative
in \hat{E} .

End of Part I

II. CLC distributions and entropy

- Approximation of $H(\mu)$ by $\sum_i H(\mu_i)$
- $p \in P_M$ as a limit of $\lambda * \mu$

Approximation of $H(\mu)$ by $\sum_i H(\mu_i)$

- Let $\mu: 2^{[n]} \rightarrow \mathbb{R}_{\geq 0}$
- Subadditivity of entropy $\rightarrow H(\mu) \leq \sum_i H(\mu_i)$
- **Theorem 5.2.** If μ is *log-concave*, then $H(\mu) \geq \sum_i \mu_i \log \left(\frac{1}{\mu_i} \right)$.

$$\mu: \sum_{S \subseteq [n]} \mu_S$$

Whiteboard

- **Theorem 5.2.** If μ is *log-concave*, then $H(\mu) \geq \sum_i \mu_i \log \left(\frac{1}{\mu_i} \right)$.

$$g_\mu(z) := \sum_s \mu_s z^s$$

$$h(z_1, \dots, z_n) := \log g_\mu \left(\frac{z_1}{\mu_1}, \dots, \frac{z_n}{\mu_n} \right).$$

Then h is a concave function.

$z \sim \mathcal{I}_s$ according to μ .

$$\text{Jensen} \Rightarrow \mathbb{E}(h(z)) \leq h(\mathbb{E}(z))$$

$$\text{RHS: } h(\mathbb{E}(z)) = h(\mu_1, \dots, \mu_n)$$

$$= \log g_\mu(1, 1, \dots, 1) = \log \left(\sum_s \mu_s \right) = 0.$$

Whiteboard

- **Theorem 5.2.** If μ is *log-concave*, then $H(\mu) \geq \sum_i \mu_i \log \left(\frac{1}{\mu_i} \right)$.

$$h(z_1, \dots, z_n) = \log g_{\mu} \left(\frac{z_1}{\mu_1}, \dots, \frac{z_n}{\mu_n} \right)$$

$$\text{LHS} = \mathbb{E}(h(z))$$

$$= \sum_S \mu_S \cdot h(\mathbb{1}_S)$$

$$= \sum_S \mu_S \cdot \log \left(\sum_T \mu_T \cdot \left(\frac{z^T}{\mu} \right) \Big| \mathbb{1}_S \right)$$

$$\geq \sum_S \mu_S \cdot \log \left(\mu_S \cdot \prod_{i \in S} \frac{1}{\mu_i} \right)$$

$$\begin{aligned} &= \underbrace{\sum_S \mu_S \cdot \log \mu_S}_{-H(\mu)} + \underbrace{\sum_S \mu_S \cdot \sum_{i \in S} \log \left(\frac{1}{\mu_i} \right)}_{= \sum_i \log \left(\frac{1}{\mu_i} \right) \cdot \sum_{S \ni i} \mu_S} \end{aligned}$$

Consequences

- **Corollary 5.7.**

$\max\left(\frac{1}{2}\sum_i H(\mu_i), \sum_i H(\mu_i) - r\right) \leq H(\mu) \leq \sum_i H(\mu_i)$ for μ : uniform distribution on bases of M

$$\begin{aligned}
 0 = \text{RHS} \geq \text{LHS} &\geq -H(\mu) + \sum_i \mu_i \log\left(\frac{1}{\mu_i}\right) \\
 \Rightarrow H(\mu) &\geq \sum_i \mu_i \log\left(\frac{1}{\mu_i}\right) \\
 \hline
 H(\mu) &\geq \sum_i \left[\mu_i \log\left(\frac{1}{\mu_i}\right) + (1-\mu_i) \log\left(\frac{1}{1-\mu_i}\right) \right] \\
 &\quad - \sum_i (1-\mu_i) \log\left(\frac{1}{\mu_i}\right) \\
 &\geq \sum_i H(\mu_i) - \sum_i \mu_i \quad \left(x \geq (1-x) \log\left(\frac{1}{1-x}\right) \right. \\
 &= \sum_i H(\mu_i) - r \quad \left. \text{for } x \in [0,1] \right)
 \end{aligned}$$

μ^* : uniform dist. on bases of M^* .
 $\mu_i^* = 1 - \mu_i$
 $H(\mu) + H(\mu^*)$
 $\geq \sum_i \mu_i \log\left(\frac{1}{\mu_i}\right) + \sum_i (1-\mu_i) \log\left(\frac{1}{1-\mu_i}\right)$
 $= \sum_i H(\mu_i)$

$p \in P_M$ as a limit of $\lambda * \mu$

- P_M : matroid polytope, i.e. $P_M = \text{conv}(1_B : B \in B_M)$

Lemma 5.8. *Let M be a matroid on ground set $[n]$ and let p be a point in \mathcal{P}_M . Then there is a distribution $\tilde{\mu}$ supported on \mathcal{B}_M with marginals p , i.e., $\tilde{\mu}_i = p_i$, such that both $\tilde{\mu}$ and $\tilde{\mu}^*$ are completely log-concave. Furthermore $\tilde{\mu}$ and $\tilde{\mu}^*$ can be obtained as the limit of external fields applied to μ and μ^* , where μ is the uniform distribution on \mathcal{B}_M .*

End of Part II

III. Max. entropy programs and basis counting

- Counting $|B_M|$
- Counting $|B_M \cap B_N|$

Counting $|B_M|$

- Let $\tau = \max(\sum_i H(p_i) \mid p = (p_1, \dots, p_n) \in P_M)$.
- To estimate $|B_M|$, output $\beta = e^\tau$.
- Then, β satisfies the following estimation guarantee:

$$\max(\sqrt{\beta}, e^{-r} \beta) \leq |B_M| \leq \beta$$

- This is the same as saying

$$\max\left(\frac{\tau}{2}, \tau - r\right) \leq \log(|B_M|) \leq \tau$$

Whiteboard

- Let $\tau = \max(\sum_i H(p_i) \mid p = (p_1, \dots, p_n) \in P_M)$.
Then, $\max\left(\frac{\tau}{2}, \tau - r\right) \leq \log(|B_M|) \leq \tau$

- μ : uniform dist. on B_M , then $\mu \in P_M$, and so

$$\log |B_M| = H(\mu) \leq \sum_i H(\mu_i) \leq \tau$$

- Given p : optimal, can find $\tilde{\mu}$: CLC, s.t. $\tilde{\mu}_i = p_i$.

$$H(\mu) \geq H(\tilde{\mu}) \geq \max\left(\frac{1}{2} \sum_{i=1}^n H(\tilde{\mu}_i), \sum_{i=1}^n H(\tilde{\mu}_i) - r\right)$$

$\tilde{\mu}$ supp. on B_M .

$$\geq \max\left(\frac{\tau}{2}, \tau - r\right).$$

Counting $|B_M \cap B_N|$ *rank-r.*

- Let $\tau = \max(\sum_i H(p_i) \mid p = (p_1, \dots, p_n) \in P_M \cap P_N)$.
- To estimate $|B_M \cap B_N|$, output $\beta = e^\tau$.
- Then, β satisfies the following estimation guarantee:

$$e^{-O(r)} \beta \leq |B_M \cap B_N| \leq \beta$$

- This is the same as saying

$$\tau - O(r) \leq \log(|B_M \cap B_N|) \leq \tau$$

Setup

$$g_M(y) \cdot g_{N^*}(z)$$

- Consider the polynomial $g(y, z) := g_M(y) \cdot g_{N^*}(z)$
- Then, $(\prod_i (\partial y_i + \partial z_i)) g(y, z) = |B_M \cap B_N|$

$$\sum_S (\partial y_S) (\partial z_{[n] \setminus S}) g_M(y) \cdot g_{N^*}(z)$$

$$T_y \supseteq T_z$$

$$S \subseteq T_y, [n] \setminus S \subseteq T_z$$

$$S \in B_M, [n] \setminus S \in B_{N^*}$$

$$S \in B_M \cap B_N.$$

- Target is to *lower-bound* LHS

Statement of the main theorem

Theorem 7.1. Let $g \in \mathbb{R}[y_1, \dots, y_n, z_1, \dots, z_n]$ be a completely log-concave multiaffine polynomial and $p \in [0, 1]^n$. Then,

$$\left(\prod_{i=1}^n (\partial_{y_i} + \partial_{z_i}) \right) g(y, z) \Big|_{y=z=0} \geq \phi(p) \cdot \inf_{y, z \in \mathbb{R}_{>0}^n} \frac{g(y, z)}{y^p z^{1-p}},$$

where $\phi(p)$ is independent of the polynomial g , and is given by the following expression:

$$\phi(p) = \prod_{i=1}^n \left(p_i^{p_i} \cdot (1 - p_i)^{1-p_i} \cdot \frac{1}{1 + p_i(1 - p_i)} \right).$$

Implication

Theorem 7.1. Let $g \in \mathbb{R}[y_1, \dots, y_n, z_1, \dots, z_n]$ be a completely log-concave multi-affine polynomial and $p \in [0, 1]^n$. Then,

$$\left(\prod_{i=1}^n (\partial_{y_i} + \partial_{z_i}) \right) g(y, z) \Big|_{y=z=0} \geq \phi(p) \cdot \inf_{y, z \in \mathbb{R}_{>0}^n} \frac{g(y, z)}{y^p z^{1-p}},$$

where $\phi(p)$ is independent of the polynomial g , and is given by the following expression:

$$\phi(p) = \prod_{i=1}^n \left(p_i^{p_i} \cdot (1-p_i)^{1-p_i} \cdot \frac{1}{1+p_i(1-p_i)} \right).$$

Take log, plug in $g = g_M \cdot g_{N^*}$,

$$\log \left(\inf_{y > 0} \frac{g_M(y)}{y^p} \right) + \log \left(\inf_{z > 0} \frac{g_{N^*}(z)}{z^{1-p}} \right)$$

When $p \in P_M \cap P_N$, find v, w , $v_i = p_i$, $w_i = 1-p_i$.
coeff of \sum in g_M .

$$\log \left(\inf_{y > 0} \frac{g_M(y)}{y^p} \right) = \sum_S v_S \log \left(\frac{c_S}{v_S} \right) = \sum_{S \in B_M} v_S \log \left(\frac{1}{v_S} \right) = H(v)$$

$$H(v) \geq \sum_i v_i \log \left(\frac{1}{v_i} \right) = \sum_i p_i \log \left(\frac{1}{p_i} \right)$$

$$\log \left(\inf_{z > 0} \frac{g_{N^*}(z)}{z^{1-p}} \right) = H(w) = H(w^*) \geq \sum_i p_i \log \left(\frac{1}{p_i} \right)$$

Implication (cont'd)

$$\left(\phi(p) = \prod_i \frac{p_i^{p_i} (1-p_i)^{1-p_i}}{1+p_i(1-p_i)} \right)$$

$$\log |B_{u_1} \cap B_{u_2}| \geq \log \phi(p) + 2 \sum_i p_i \log \left(\frac{1}{p_i} \right)$$

$$\left(\frac{p^p (1-p)^{1-p}}{1+p(1-p)} \geq p \log \left(\frac{1}{e^2} \right) \right)$$

$$\geq \sum_i p_i \log \left(\frac{1}{e^2} \right) + 2 \sum_i p_i \log \left(\frac{1}{p_i} \right)$$

$$= \sum_i p_i \log \left(\frac{1}{p_i} \right) - 2 \sum_i p_i$$

$$= \sum_i H(p_i) - 2r - \sum_i (1-p_i) \cdot \log \frac{1}{1-p_i}$$

$$\geq \sum_i H(p_i) - 3r.$$

Base case ($n = 1$)

- Characterize log-concavity of $g(y, z) = a + by + cz + dyz$
 - Assume $a, b, c, d \geq 0$

- Prove that

$$(\partial y + \partial z)g(y, z) \Big|_{y=z=0} \geq \frac{p^p(1-p)^{1-p}}{1+p(1-p)} \inf_{y,z>0} \frac{g(y, z)}{y^p z^{1-p}}$$

Whiteboard

- Characterize log-concavity of $g(y, z) = a + by + cz + dyz$

$$g(y, z) \text{ log-concave} \Leftrightarrow 2bc - ad \geq 0.$$

Whiteboard

- Prove that

$$(\partial_y + \partial_z)g(y, z) \Big|_{y=z=0} \geq \frac{p^p(1-p)^{1-p}}{1+p(1-p)} \inf_{y, z > 0} \frac{g(y, z)}{y^p z^{1-p}}$$

$$g(y, z) = a + by + cz + dyz$$

$$\text{LHS} = b + c$$

Look at $\inf_{y, z > 0} \frac{g(y, z)}{y^p z^{1-p}}$

- Special case 1: $p \in (0, 1)$

Say $p=0$: $\inf_{y, z > 0} \frac{a + by + cz + dyz}{z}$

$$\begin{aligned} y &\rightarrow 0 \\ z &\rightarrow +\infty \\ &\searrow c \end{aligned}$$

- Special case 2: $d=0$.

$\hookrightarrow bc=0$, $b=0$, let $y \rightarrow +\infty$, $\inf = 0$.

$\hookrightarrow bc > 0$, by $zbc \geq ad$, scale $d \uparrow$

Whiteboard

- Prove that

$$(\partial_y + \partial_z)g(y, z) \Big|_{y=z=0} \geq \frac{p^p(1-p)^{1-p}}{1+p(1-p)} \inf_{y, z > 0} \frac{g(y, z)}{y^p z^{1-p}}$$

$2bc = ad$, Scale so that $d=1$

$$\inf_{y, z > 0} \frac{2bc + by + cz + yz}{y^p z^{1-p}} = \inf_{y, z > 0} \frac{bc + (c+y)(b+z)}{y^p z^{1-p}}$$

$$\rightarrow y = c \cdot \frac{p}{1-p}, \quad z = b \cdot \frac{1-p}{p}$$

$$\left(y = c \cdot \left[\frac{p}{1-p} + \sqrt{1 + \left(\frac{p}{1-p}\right)^2} - 1 \right], \quad z = b \cdot \left[\frac{1-p}{p} + \sqrt{1 + \left(\frac{1-p}{p}\right)^2} - 1 \right] \right)$$

$p = \frac{1}{2}, 5 \quad \text{vs} \quad (2+2\sqrt{2})$

Whiteboard

- Prove that

$$(\partial y + \partial z)g(y, z) \Big|_{y=z=0} \geq \frac{p^p(1-p)^{1-p}}{1+p(1-p)} \inf_{y, z > 0} \frac{g(y, z)}{y^p z^{1-p}}$$

$$\text{Plus } y = c \frac{p}{1-p}, \quad z = b \frac{1-p}{p}.$$

$$\text{RHS} \leq \frac{p^p (1-p)^{1-p}}{1+p(1-p)} \cdot \frac{bc \left[1 + \frac{1}{1-p} \cdot \frac{1}{p}\right]}{c^p b^{1-p} \left(\frac{p}{1-p}\right)^p \cdot \left(\frac{1-p}{p}\right)^{1-p}}$$

$$= b^p c^{1-p} \cdot \frac{1+p(1-p)}{p(1-p)} \cdot (1-p)^p \cdot p^{1-p} \cdot \frac{1}{1+p(1-p)}$$

$$= \left(\frac{b}{p}\right)^p \cdot \left(\frac{c}{1-p}\right)^{1-p}$$

$$\leq b + c = \text{LHS}.$$

End of Part III